Lecture #8: UTXOs and Accounts

COMS 4995-001: The Science of Blockchains URL: https://timroughgarden.org/s25/

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Responsibilities of a Blockchain Protocol

Consensus: decide on a sequence (aka "chain") of blocks.

- note: all validators must agree on this sequence!
- blocks keeping getting added (one-by-one) as long as there are transactions to process
- SMR, Tendermint vs. longest-chain consensus, etc.

The Consensus Layer



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Execution: keep state of the virtual machine up-to-date.

- new block added → execute the corresponding snippets of code (do computations, update variable values, etc.)
- subject of this week (concludes Part I of course)

The Computer in the Sky



The Execution Layer

tx2 tx3	tx1	tx4
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consensus transaction sequence

The Execution Layer



consensus transaction sequence



blockchain protocol (execution layer) [replicated at each physical machine]

The Execution Layer



simulated (virtual) computer [replicated at each physical machine]

[replicated at each physical machine]

Recap: A Cartoon of Web3



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Blockchain protocol:

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 - acts as a "master program" to coordinate all apps/smart contracts
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 - provides a virtual machine to developers of applications
- like the Internet, "decentralized" -- the product of collaboration between many physical machines, no one owner/operator

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 - acts as a "master program" to coordinate all apps/smart contracts
 - provides a virtual machine to developers of applications
- "decentralized" like the Internet

Goals for Lecture #8

- 1. The UTXO model (used, e.g., in Bitcoin).
 - counterintuitive but elegant VM specialized for payments
- 2. Measuring the size of a transaction.
 - idea: what resources are required (now and forever) by a transaction?
 - in practice, very tricky!
- 3. The account-based model (used in Ethereum and Solana).
 - explicit notion of account IDs and balances, programs as accounts
- 4. Metering computation.

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Next: warm-up with deep dive on Bitcoin's "execution layer."

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input 2	tx ID = 1713	output 2
		output 3

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- format for one input
 - output of some other tx [should be unspent! → "UTXO"]
 - "witness" satisfying output's spending condition [typically, ≈ a signature]



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- for each input, witness satisfies UTXO's spending conditions



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• canonical implementation = "Bitcoin core"

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Challenge: multiple types of resources required:

- resources at the consensus layer (bandwidth)
- resources at the execution layer (computation, memory access)
- resources for long-term storage (at validators or elsewhere)

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Further challenge: resource consumption may depend on external-to-protocol factors (e.g., specific client implementation and/or validator architecture).

- in practice, "size" often defined w.r.t. some canonical implementation

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- lowers barrier to participating, but raises barrier to validating
- benefits of innovating vs. benefits of hardening
- led to Bitcoin Cash (fork of Bitcoin with bigger blocks, now irrelevant)

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- idea: validator can discard witnesses after checking tx validity
 - "archival nodes" should still keep witness data for posterity

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- debate: are these good for Bitcoin?

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Point: definition of transaction size can fundamentally affect how a blockchain protocol is used!

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Example: in Ethereum, a user account ("EOA") has no code, only data is a "nonce" [= # of txs sent by account, prevents "replay attacks"].

• all other data on user stored in contracts' accounts

- signature by the sender (can back out pk/ID from signature)
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- declaration of resources to be used
- transaction fee

Transaction: sent by a user (to a user, or a program). Includes:

- signature by the sender (can back out pk/ID from signature)
- recipient (specified by account ID, user or contract)
- value (in native currency)
- data (e.g., which function to call and with which arguments)
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Note: if programs can be arbitrary code, corresponding state transition can be extremely complex.

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- associate an amount of "gas" with each EVM opcode
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 - add two numbers = 3 units of gas; evaluate SHA-256 = 30 units
- user prepays for gas (part of the tx description)
- run out of gas mid-execution
 → tx aborted and rolled back 69